	and <b>errno</b> is set to indicate the error.	<b>RETURN VALUES</b> If a function in the exec family returns to the calling process, an error has occurred; the return value is $-1$	process.	signals that are being caugin by the caning process are set to the becaut disposition in the new process image (see signal(3C)). Otherwise, the new process image inherits the signal dispositions of the calling	File descriptors open in the calling process remain open in the new process.	The <i>file</i> argument points to the new process file. If <i>file</i> does not contain a slash character, the path prefix for this file is obtained by a search of the directories passed in the <b>PATH</b> environment variable (see <b>environ</b> ( $5$ )).	The <i>path</i> argument points to a path name that identifies the new process file.	by a null pointer.	it should point to a string that is the same as <i>path</i> (or its last component). The <i>argv</i> argument is terminated	argument list available to the new process image. By convention, $ar_{gv}$ must have at least one member, and	The <i>argy</i> argument is an array of character pointers to null-terminated strings. These strings constitute the	strings is terminated by a (chars)0 argument.	ment is a variatione to une new process intages. Conventionality at reast $a_{ij}$ or should be present. The $a_{ij}$ or a variance that is the same as <i>rath</i> (or the last commonant to $f$ <i>arch</i> ). The list of a ranneat	The arguments $argo, \ldots, argo point to null-terminated character strings. These strings constitute the arguments that are the arguments arguments that the string s$	emp is an analy or character pointers to the environment sumps. As indicated, $a/g$ is at least one, and the first member of the stray notions to a strip containing the name of the file	where <i>argc</i> is the argument count, <i>argv</i> is an array of character pointers to the arguments themselves, and	листиали (лисатусу спол "от. 5 г.), спол "стиурДУ)	int main (int area char *arov) char *envn().	When a C program is executed, it is called as follows:	the calling process image is overlaid by the new process image.	of data for an interpreter. There can be no return from a successful call to one of these functions because	Each of the functions in the <b>exec</b> family overlays a new process image on an old process. The new process image is constructed from an ordinary executable file. This file is either an executable object file or a file	DESCRIPTION	int execvp (const char * <i>fle</i> , char *const $argv[I)$ ;	int execlp (const char $*$ <i>file</i> , const char $*arg\theta$ ,, const char $*argn$ , char $*/*$ NULL*/);	int execve (const char * <i>path</i> , char *const <i>argv[ ]</i> char *const <i>emp[ ]</i> );	char */*NULL*, char *const emp[]);	int execle(const char $path$ , char $const arg0[1, \dots, const char argn,$	int execv(const char * <i>path</i> , char *const <i>argv[ ]</i> );	int execl(const char * <i>path</i> , const char * <i>arg0</i> ,, const char * <i>argn</i> , char * /*NULL*/);	#include <unistd.h></unistd.h>	SYNOPSIS	exec, execl, execv, execle, execve, execlp, execvp – execute a file	NAME
ENVENT A component of <i>pain</i> uses not exist, or <i>pain</i> is an empty string.		<b>RETURN VALUE</b> On success, zero is returned. On error, -1 is returned, and <i>errno</i> is set appropriately.	S_ISDIR(m) directory?	_ISREG(m)	The following POSIX macros are defined to check the file type in the field st_mode:	is the length of the pathname it contains, without trailing NUL.	). The value <i>st size</i> gives the size of the file (if it is a regular file or a symlink) in bytes. The size of a symlink	t. time_t st_cume; /* time of fast status change */	st_mtime;	st_atime;		Neize t «t Nlečize /* Holokijze for filesvitem I/O */	st_rdev;	st_gid;	 st_nlink; /	mode_t st_mode; /* protection */	st_ucv,		struct stat (	They all return a <i>stat</i> structure, which contains the following fields:		<b>istat</b> is identical to <b>stat</b> , except in the case of a symbolic link, where the link itself is stat-ed, not the file that it refers to		<b>stat</b> stats the file pointed to by <i>path</i> and fills in <i>buf</i> .	get this information but you need search rights to all directories named in the path leading to the file.	DESCRIPTION These functions return information about the specified file. You do not need any access rights to the file to	in san(const that $r$ join, struct stat $roig$ );	int stat(const char * path, struct stat * buf);		#include <sys stat.n=""> #include <unistd.h></unistd.h></sys>	#include <sys types.h=""></sys>	SYNOPSIS	stat, lstat – get file status	NAME

-	H	H	ORS	
ENOTDIR	ENOENT	EACCES		
<b>DIR</b> A component of the path prefix of <i>path</i> is not a directory.	2NT A component of <i>path</i> does not exist, or <i>path</i> is an empty string.	<b>ES</b> Search permisson is denied for one of the directories in the path prefix of <i>path</i> .		

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exec(2)

exec(2)

stat(2)

wait(2)

wait(2)

NAME

wait, waitpid, waitid - wait for process to change state

# SYNOPSIS

#include <sys/types.h>
#include <sys/wait.h>

pid\_t wait(int \*status);

pid\_t waitpid(pid\_t pid, int \*status, int options);

#### DESCRIPTION

All of these system calls are used to wait for state changes in a child of the calling process, and obtain information about the child whose state has changed. A state change is considered to be: the child terminated; the child was stopped by a signal; or the child was resumed by a signal. In the case of a terminated child, performing a wait allows the system to release the resources associated with the child; if a wait is not performed, then the terminated child remains in a "zombie" state (see NOTES below).

If a child has already changed state, then these calls return immediately

**waitpid**() system call suspends execution of the calling process until a child specified by *pid* argument has changed state. By default, **waitpid**() waits only for terminated children, but this behavior is modifiable via the options argument, as described below. The wait() system call suspends execution of the calling process until one of its children terminates. The

The value of pid can be:

- Â meaning wait for any child process whose process group ID is equal to the absolute value of pid.
- Ļ meaning wait for any child process
- 0 meaning wait for any child process whose process group ID is equal to that of the calling process
- $\stackrel{\vee}{0}$ meaning wait for the child whose process ID is equal to the value of pid.

The value of options is an OR of zero or more of the following constants:

WNOHANG return immediately if no child has exited.

to it, as is done in **wait()** and **waitpid()**!): ger can be inspected with the following macros (which take the integer itself as an argument, not a pointer If status is not NULL, wait() and waitpid() store status information in the int to which it points. This inte-

## WIFEXITED(status)

from main(). returns true if the child terminated normally, that is, by calling exit(3) or \_exit(2), or by returning

WEXITSTATUS(status)

ment in main(). This macro should only be employed if WIFEXITED returned true. ment that the child specified in a call to exit(3) or  $_exit(2)$  or as the argument for a return statereturns the exit status of the child. This consists of the least significant 8 bits of the status argu-

### RETURN VALUE

wait(): on success, returns the process ID of the terminated child; on error, -1 is returned

waitpid(): on success, returns the process ID of the child whose state has changed; if WNOHANG was returned. On error, -1 is returned. specified and one or more child(ren) specified by *pid* exist, but have not yet changed state, then 0 is