accept(2)

accept(2)

bind(2)

bind(2)

NAME

accept - accept a connection on a socket

SYNOPSIS #include <sys/types.h>

#include <sys/socket.h>

int accept(int s, struct sockaddr *addr, int *addrlen);

DESCRIPTION

more connections. The original socket (s) remains open for accepting further connections. socket, ns, is used to read and write data to and from the socket that connected to ns; it is not used to accept name associated with s. This is the device on which the connect indication will be accepted. The accepted described below. The accept() function uses the netconfig(4) file to determine the STREAMS device file marked as non-blocking and no pending connections are present on the queue, accept() returns an error as not marked as non-blocking, accept() blocks the caller until a connection is present. If the socket is new file descriptor, ns, for the socket. If no pending connections are present on the queue and the socket is nection on the queue of pending connections, creates a new socket with the properties of s, and allocates a The argument s is a socket that has been created with **socket**(3N) and bound to an address with **bind**(3N), and that is listening for connections after a call to **listen**(3N). The **accept**() function extracts the first con-

ERRORS

the global errno

EACCES

to access it.

The requested address is protected and the current user has inadequate permission

The **bind()** call will fail if:

RETURN VALUES

assigned to the socket.

space (address family) but has no name assigned. bind() requests that the name pointed to by name be

bind() assigns a name to an unnamed socket. When a socket is created with socket(3N), it exists in a name

If the bind is successful, 0 is returned. A return value of -1 indicates an error, which is further specified in

DESCRIPTION

int bind(int s, const struct sockaddr *name, int namelen);

#include <sys/socket.h>

#include <sys/types.h>

SYNOPSIS NAME

bind - bind a name to a socket

in which the communication occurs. known to the communications layer. The exact format of the addr parameter is determined by the domain The argument *addr* is a result parameter that is filled in with the address of the connecting entity as it is

addr; on return it contains the length in bytes of the address returned. The argument *addrlen* is a value-result parameter. Initially, it contains the amount of space pointed to by

EBADF

s is not a valid descriptor.

The specified address is not available on the local machine

The specified address is already in use.

EADDRNOTAVAIL EADDRINUSE

The accept() function is used with connection-based socket types, currently with SOCK_STREAM.

accept(). read. However, this will only indicate when a connect indication is pending; it is still necessary to call It is possible to select(3C) or poll(2) a socket for the purpose of an accept() by selecting or polling it for a

RETURN VALUES

tor for the accepted socket. The accept() function returns -1 on error. If it succeeds, it returns a non-negative integer that is a descrip-

ERRORS

| ac | |
|----------|--|
| accept() | |
| will f | |
| fail if: | |
| | |

| beel | EPROTO A p | ENOMEM The | ENODEV The pro- fig file. | EMFILE The | EINTR The | EBADF The | ассери() win тап н: |
|---|--|---|--|---|---|----------------------------|---------------------|
| The socket is marked as non-blocking and no connections are present to be | A protocol error has occurred; for example, the STREAMS protocol stack has not been initialized or the connection has already been released. | There was insufficient user memory available to complete the operation. | The protocol family and type corresponding to <i>s</i> could not be found in the netcon-fig file. | The per-process descriptor table is full. | The accept attempt was interrupted by the delivery of a signal. | The descriptor is invalid. | |

NOTES SEE ALSO

unlink(2), socket(3N), attributes(5), socket(5)

EROFS

The inode would reside on a read-only file system.

A component of the path prefix of the pathname in *name* is not a directory. A component of the path prefix of the pathname in *name* does not exist. Too many symbolic links were encountered in translating the pathname in name.

ENOTDIR ENOENT ELOOP EISDIR EIO

A null pathname was specified

An I/O error occurred while making the directory entry or allocating the inode.

Search permission is denied for a component of the path prefix of the pathname in

EACCES

name.

ENOTSOCK ENOSR EINVAL EINVAL

The following errors are specific to binding names in the UNIX domain

s is a descriptor for a file, not a socket.

There were insufficient STREAMS resources for the operation to complete

The socket is already bound to an address

namelen is not the size of a valid address for the specified address family

Binding a name in the UNIX domain creates a socket in the file system that must be deleted by the caller

The rules used in name binding vary between communication domains

when it is no longer needed (using unlink(2)).

SEE ALSO

poll(2), bind(3N), connect(3N), listen(3N), select(3C), socket(3N), netconfig(4), attributes(5), socket(5)

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dup(2)

dup(2)

exec(2)

NAME

SYNOPSIS dup, dup2 - duplicate a file descriptor

#include <unistd.h>

int dup2(int oldfd, int newfd); int dup(int oldfd);

DESCRIPTION

dup() and dup2() create a copy of the file descriptor *oldfd*

dup() uses the lowest-numbered unused descriptor for the new descriptor.

dup2() makes *newfd* be the copy of *oldfd*, closing *newfd* first if necessary, but note the following:

- If *oldfd* is not a valid file descriptor, then the call fails, and *newfd* is not closed.
- If *oldfd* is a valid file descriptor, and *newfd* has the same value as *oldfd*, then **dup2()** does nothing, and returns newfa

changed for the other. flags; for example, if the file offset is modified by using lseek(2) on one of the descriptors, the offset is also ably. They refer to the same open file description (see open(2)) and thus share file offset and file status After a successful return from dup() or dup2(), the old and new file descriptors may be used interchange-

The two descriptors do not share file descriptor flags (the close-on-exec flag). The close-on-exec flag (FD_CLOEXEC; see **fcntl**(2)) for the duplicate descriptor is off.

RETURN VALUE

ately). dup() and dup2() return the new descriptor, or -1 if an error occurred (in which case, *errno* is set appropri-

ERRORS

EBADF

oldfd isn't an open file descriptor, or *newfd* is out of the allowed range for file descriptors.

EBUSY

(Linux only) This may be returned by dup2() during a race condition with open(2) and dup().

EINTR The **dup2**() call was interrupted by a signal; see **signal**(7).

EMFILE

The process already has the maximum number of file descriptors open and tried to open a new

CONFORMING TO

SVr4, 4.3BSD, POSIX. 1-2001

NOTES

of range. On some systems dup2() also sometimes returns EINVAL like F_DUPFD. The error returned by dup2() is different from that returned by fcnt1(..., F_DUPFD, ...) when newfd is out

mer will not use dup2() without closing newfd first. If newfd was open, any errors that would have been reported at close(2) time are lost. A careful program-

SEE ALSO

close(2), fcntl(2), open(2)

COLOPHON This page is part of release 3.05 of the Linux man-pages project. A description of the project, and informa-

tion about reporting bugs, can be found at http://www.kemel.org/doc/man-pages/

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NAME

exec, execl, execv, execle, execve, execlp, execvp - execute a file

SYNOPSIS

#include <unistd.h>

int execv(const char * path, char *const argv[]); int execl(const char * path, const char * arg0, ..., const char * argn, char * /*NULL*/);

int execle(const char * path, char * const arg0[], ..., const char * argn,

char * /*NULL*/, char *const envp[]);

int execve (const char * path, char * const argv[] char * const envp[]);

int execlp (const char * *file*, const char * *arg0*, ..., const char * *argn*, char * /*NULL*/);

int execvp (const char * file, char * const argv[]);

DESCRIPTION

Each of the functions in the **exec** family overlays a new process image on an old process. The new process image is constructed from an ordinary, executable file. This file is either an executable object file, or a file of data for an interpreter. There can be no return from a successful call to one of these functions because the calling process image is overlaid by the new process image.

When a C program is executed, it is called as follows:

int main (int argc, char *argv[], char *envp[]);

first member of the array points to a string containing the name of the file. envp is an array of character pointers to the environment strings. As indicated, argc is at least one, and the where argc is the argument count, argv is an array of character pointers to the arguments themselves, and

ment list available to the new process image. Conventionally at least arg0 should be present. The arg0strings is terminated by a (char *)0 argument argument points to a string that is the same as *path* (or the last component of *path*). The list of argument The arguments arg0, ..., argn point to null-terminated character strings. These strings constitute the argu-

argument list available to the new process image. By convention, *argv* must have at least one member, and it should point to a string that is the same as *path* (or its last component). The *argv* argument is terminated by a null pointer. The argv argument is an array of character pointers to null-terminated strings. These strings constitute the

The path argument points to a path name that identifies the new process file.

The file argument points to the new process file. If file does not contain a slash character, the path prefix for this file is obtained by a search of the directories passed in the **PATH** environment variable (see **environ**(5)).

File descriptors open in the calling process remain open in the new process.

process. image (see signal(3C)). Otherwise, the new process image inherits the signal dispositions of the calling Signals that are being caught by the calling process are set to the default disposition in the new process

RETURN VALUES

and errno is set to indicate the error. If a function in the exec family returns to the calling process, an error has occurred; the return value is -1

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fopen/fdopen/fileno(3)

fopen/fdopen/fileno(3)

NAME

fopen, fdopen, fileno - stream open functions

SYNOPSIS #include <stdio.h>

int fileno(FILE *stream); FILE *fdopen(int fildes, const char *mode); FILE *fopen(const char * path, const char *mode);

DESCRIPTION

The **fopen** function opens the file whose name is the string pointed to by *path* and associates a stream with

may follow these sequences.): The argument mode points to a string beginning with one of the following sequences (Additional characters

- ٦ Open text file for reading. The stream is positioned at the beginning of the file
- 7 Open for reading and writing. The stream is positioned at the beginning of the file.
- ¥ Truncate file to zero length or create text file for writing. The stream is positioned at the beginning
- \mathbf{W}_{+} of the file. stream is positioned at the beginning of the file. Open for reading and writing. The file is created if it does not exist, otherwise it is truncated. The
- a positioned at the end of the file. Open for appending (writing at end of file). The file is created if it does not exist. The stream is
- a+ The stream is positioned at the end of the file. Open for reading and appending (writing at end of file). The file is created if it does not exist.

dup'ed, and will be closed when the stream created by fdopen is closed. The result of applying fdopen to a The **fdopen** function associates a stream with the existing file descriptor, *fildes*. The *mode* of the stream (one of the values "r", "r+", "w", "w+", "a", "a+") must be compatible with the mode of the file descriptor. shared memory object is undefined. indicators are cleared. Modes "w" or "w+" do not cause truncation of the file. The file descriptor is not The file position indicator of the new stream is set to that belonging to *fildes*, and the error and end-of-file

The function fileno() examines the argument stream and returns its integer descriptor.

RETURN VALUE

Upon successful completion **fopen**, **fdopen** and **freopen** return a **FILE** pointer. Otherwise, **NULL** is returned and the global variable *errno* is set to indicate the error.

ERRORS

EINVAL

The mode provided to fopen, fdopen, or freopen was invalid.

routine malloc(3). The **fopen**, **fdopen** and **freopen** functions may also fail and set *errno* for any of the errors specified for the

The **fopen** function may also fail and set *errno* for any of the errors specified for the routine **open**(2).

The **fdopen** function may also fail and set errno for any of the errors specified for the routine **fcntl**(2)

SEE ALSO open(2), fclose(3), fileno(3)

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gets(3)

NAME

fputs, puts - output of strings gets, fgets - get a string from a stream

SYNOPSIS

#include <stdio.h>

char *gets(char *s);

char *fgets(char *s, int n, FILE *stream);

int fputs(const char *s, FILE *stream);

int puts(const char *s);

DESCRIPTION gets/fgets

character is discarded and the string is terminated with a null character. The gets() function reads characters from the standard input stream (see intro(3)), stdin, into the array pointed to by s, until a newline character is read or an end-of-file condition is encountered. The newline

are read, or a newline character is read and transferred to s, or an end-of-file condition is encountered. The string is then terminated with a null character. The fgets() function reads characters from the stream into the array pointed to by s, until n-1 characters

For this reason, it is strongly recommended that gets() be avoided in favor of fgets(). When using gets(), if the length of an input line exceeds the size of s, indeterminate behavior may result.

RETURN VALUES

If end-of-file is encountered and no characters have been read, no characters are transferred to s and a null encountered, the EOF indicator for the stream is set. Otherwise s is returned. opened for reading, a null pointer is returned and the error indicator for the stream is set. If end-of-file is pointer is returned. If a read error occurs, such as trying to use these functions on a file that has not been

ERRORS

The gets() and fgets() functions will fail if data needs to be read and:

EOVERFLOW The file is a regular file and an attempt was made to read at or beyond the offset maximum associated with the corresponding stream.

DESCRIPTION puts/fputs

fputs() writes the string *s* to *stream*, without its trailing '\0'

puts() writes the string s and a trailing newline to stdout

from the stdio library for the same output stream. Calls to the functions described here can be mixed with each other and with calls to other output functions

RETURN VALUE

puts() and fputs() return a non - negative number on success, or EOF on error

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socket(2) / ipv6(7)

socket(2) / ipv6(7)

listen(2)

SYNOPSIS NAME

listen - listen for connections on a socket

listen(2)

NAME

ipv6, PF_INET6 - Linux IPv6 protocol implementation

SYNOPSIS #include <sys/socket.h>

#include <netinet/in.h>

udp6_socket = socket(PF_INET6, SOCK_DGRAM, protocol); raw6_socket = socket(PF_INET6, SOCK_RAW, protocol); tcp6_socket = socket(PF_INET6, SOCK_STREAM, 0);

DESCRIPTION

accept incoming connection requests using accept(2).

listen() marks the socket referred to by sockfd as a passive socket, that is, as a socket that will be used to

The sockfd argument is a file descriptor that refers to a socket of type SOCK_STREAM or SOCK_SEQ-

PACKET.

int listen(int sockfd, int backlog);

#include <sys/socket.h>

#include <sys/types.h>

/* See NOTES */

DESCRIPTION

Linux 2.2 optionally implements the Internet Protocol, version 6. This man page contains a description of the IPv6 basic API as implemented by the Linux kernel and glibc 2.1. The interface is based on the BSD sockets interface; see socket(7).

man page The IPv6 API aims to be mostly compatible with the ip(7) v4 API. Only differences are described in this

expands to a constant expression. Both of them are in network order able which has in6_addr type. In static initializations IN6ADDR_ANY_INIT may also be used, which To bind an AF_INET6 socket to any process the local address should be copied from the in6addr_any vari-

The IPv6 loopback address (::1) is available in the global in6addr_loopback variable. For initializations IN6ADDR_LOOPBACK_INIT should be used.

ERRORS

EADDRINUSE

EBADF

The argument sockfd is not a valid descriptor.

Another socket is already listening on the same port

ENOTSOCK

The argument sockfd is not a socket.

RETURN VALUE

ignored so that a later reattempt at connection succeeds.

On success, zero is returned. On error, -1 is returned, and errno is set appropriately.

The *backlog* argument defines the maximum length to which the queue of pending connections for *sockfd* may grow. If a connection request arrives when the queue is full, the client may receive an error with an indication of **ECONNREFUSED** or, if the underlying protocol supports retransmission, the request may be

the address handling functions in libc. gram only needs only to support this API type to support both protocols. This is handled transparently by IPv4 connections can be handled with the v6 API by using the v4-mapped-on-v6 address type; thus a pro-

source address will be mapped to v6 and it will be mapped to v6. IPv4 and IPv6 share the local port space. When you get an IPv4 connection or packet to a IPv6 socket its

Address Format

struct sockaddr_in6 { uint32_t sin6_flowinfo; /* IPv6 flow information */ struct in6_addr sin6_addr; /* IPv6 address */ uint32_t uint16_t uint32_t uint16_t sin6_scope_id; /* Scope ID (new in 2.4) */ sin6_port; sin6_family; /* AF_INET6 */ /* port number */

NOTES

To accept connections, the following steps are performed:

A socket is created with socket(2)

struct in6_addr { unsigned char s6_addr[16]; /* IPv6 address */

sin6_family is always set to AF_INET6; sin6_port is the protocol port (see sin_port in ip(7)); sin6_flowinfo is the IPv6 flow identifier; sin6_addr is the 128-bit IPv6 address. sin6_scope_id is an ID of depending of on the scope of the address. It is new in Linux 2.4. Linux only supports it for link scope addresses, in that case *sin6_scope_id* contains the interface index (see **netdevice**(7))

NOTES

instead. types can be stored safely in a struct sockaddr need to be changed to use struct sockaddr_storage for that The sockaddr_in6 structure is bigger than the generic sockaddr. Programs that assume that all address

> SEE ALSO EXAMPLE

See bind(2)

accept(2), bind(2), connect(2), socket(2), socket(7)

cated to that value; the default value in this file is 128

4 <u>ω</u> 2

to it.

Connections are accepted with accept(2).

specified with listen().

If the backlog argument is greater than the value in /proc/sys/net/core/somaxconn, then it is silently trun-

A willingness to accept incoming connections and a queue limit for incoming connections are

The socket is bound to a local address using **bind**(2), so that other sockets may be **connect**(2)ed

SEE ALSO

cmsg(3), ip(7)

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|---|-------|
| SEE ALSO stgaction(2), stgsetops(3C), | SEE A |
| EINTR A signal is caught by the calling process and control is returned from the signal catching function. | |
| EFAULT set points to an illegal address. | |
| ERRORS sigsuspend() fails if either of the following is true: | ERRO |
| RETURN VALUES Since sigsuspend() suspends process execution indefinitely, there is no successful completion return value. On failure, it returns -1 and sets errno to indicate the error. | RETU |
| It is not possible to block those signals that cannot be ignored (see signal (5)); this restriction is silently imposed by the system. | |
| If the action is to terminate the process, sigsuspend() does not return. If the action is to execute a signal catching function, sigsuspend() returns after the signal catching function returns. On return, the signal mask is restored to the set that existed before the call to sigsuspend() . | |
| DESCRIPTION sigsuspend sigsuspend() replaces the caller's signal mask with the set of signals pointed to by the argument set and then suspends the caller until delivery of a signal whose action is either to execute a signal catching func- tion or to terminate the process. | DESC |
| EINVAL The value of the <i>how</i> argument is not equal to one of the defined values. | |
| EFAULT set or oset points to an illegal address. | |
| sigprocmask() fails if any of the following is true: | |
| ERRORS | ERRO |
| RETURN VALUES On success, sigprocmask() returns 0. On failure, it returns -1 and sets errno to indicate the error. | RETU |
| If sigprocmask() fails, the caller's signal mask is not changed. | |
| It is not possible to block those signals that cannot be ignored this restriction is silently imposed by the sys- tem. See sigaction (2). | |
| If there are any pending unblocked signals after the call to sigprocmask(), at least one of those signals will be delivered before the call to sigprocmask() returns. | |
| SIG_BLOCK , the set pointed to by the argument <i>set</i> is added to the current signal mask. If the value is SIG_UNBLOCK , the set pointed by the argument <i>set</i> is removed from the current signal mask. If the value is SIG_SETMASK , the current signal mask is replaced by the set pointed to by the argument <i>set</i> . If the argument <i>set</i> is not NULL, the previous mask is stored in the space pointed to by <i>set</i> . If the value of the argument <i>set</i> is NULL, the value <i>how</i> is not significant and the caller's signal mask is unchanged; thus, the call can be used to inquire about currently blocked signals. | |
| DESCRIPTION sigprocmask The sigprocmask() function is used to examine and/or change the caller's signal mask. If the value is | DESC |
| <pre>int sigsuspend(const sigset_t *set);</pre> | |
| int sigprocmask(int how, const sigset_t * set, sigset_t * aset); | |
| SYNOPSIS #include <signal.h></signal.h> | SYNO |
| AME sigprocmask – change and/or examine caller's signal mask sigsuspend – install a signal mask and suspend caller until signal | NAME |

sigsuspend/sigprocmask(2)

| SP-Klausur Manual-Auszug 2012-02-22 I | | | | | SEE ALSO sigaction(2), sigpending(2), sigprocmask(2), sigsuspend(2), attributes(5), signal(5) | EFAULT The set argument specifies an invalid address. | sightliset() will fail if the following is true: | sigaddset(), sigde SINVAL | return a value of 0. Otherwise a value of -1 is returned and errno is set to indicate the error. ERRORS | KELUKA VALUES Upon successful completion, the sigismember () function returns a value of one if the specified signal is a member of the specified set, or a value of 0 if it is not. Upon successful completion, the other functions | Any object of type <i>sigeref_t</i> must be initialized by applying either sigemptyset () of signilise () before applying any other operation. | sigismember() checks whether the signal specified by the value of signo is a member of the set pointed to by set. | sigaddset () adds the individual signal specified by the value of <i>signo</i> to the set pointed to by <i>set</i> . sigdelset () deletes the individual signal specified by the value of <i>signo</i> from the set pointed to by <i>set</i> . | sigemptyset() initializes the set pointed to by set to exclude all signals defined by the system. sigfillset() initializes the set pointed to by set to include all signals defined by the system. | These functions manipulate <i>sigset_t</i> data types, representing the set of signals supported by the implementation. | DESCRIPTION | int sigatouset(sigset_t "set, int signo); int sigdelset(sigset_t "set, int signo); int sigismember(sigset_t "set, int signo); | int signillset(sigset_t * set); | <pre>int sigemptyset(sigset_t * ser);</pre> | SYNOPSIS #include <signal.h></signal.h> | NAME sigsetops, sigemptyset, sigfillset, sigaddset, sigdelset, sigismember – manipulate sets of signals | sigsetops(3C) sigsetops(3C) |
|---------------------------------------|---|--|--|--|--|---|---|--|---|---|---|---|--|--|--|--|---|---------------------------------|---|---|--|-----------------------------|
| SP-Klausur Manual-Auszug 2012-02-22 I | SEE ALSO exec(2), exit(2), fork(2), sigaction(2), wstat(5) | EINVAL An invalid value was specified for <i>options</i> . | ECHILD The process or process group specified by <i>pid</i> does not exist or is not a child of the calling process or can never be in the states specified by <i>options</i> . | ERRORS waitpid() will fail if one or more of the following is true: | Substantiation that canning process, -1 is returned and error is set to EAVIK. It this function was invoked with WNOHANG set in <i>options</i> , it has at least one child process specified by <i>pid</i> for which status is not available, and status is not available for any process specified by <i>pid</i> , 0 is returned. Otherwise, -1 is returned, and error. | the process ID of the child process for which status is reported. If waitpid () returns device dyna to the delivery of a | RETURN VALUES If waithid () rations because the status of a shild process is available, this function returns a value equal to | WNOWAIT Keep the process whose status is returned in <i>stat_loc</i> in a waitable state. The process may be waited for again with identical results. | WNOHANG waitpid() will not suspend execution of the calling process if status is not imme- diately available for one of the child processes specified by <i>pid</i> . | WCONTINUED The status of any continued child process specified by <i>pid</i> , whose status has not been reported since it continued, is also reported to the calling process. | The <i>options</i> argument is constructed from the bitwise inclusive OR of zero or more of the following flags, defined in the header <sys wait.h=""></sys> : | If waitpid () returns because the status of a child process is available, then that status may be evaluated with the macros defined by wstat (5). If the calling process had specified a non-zero value of <i>stat_loc</i> , the status of the child process will be stored in the location pointed to by <i>stat_loc</i> . | If pid is less than $(pid_t)-1$, status is requested for any child process whose process group ID is equal to the absolute value of pid . | If <i>pid</i> is equal to (pid_t)0 status is requested for any child process whose process group ID is equal to that of the calling process. | If pid is greater than (pid_10 , it specifies the process ID of the child process for which status is requested. | If <i>pid</i> is equal to (pid_t)–1 , status is requested for any child process. | wattpid() suspends the calling process until one of its children changes state; if a child process changed state prior to the call to waitpid() , return is immediate. <i>pid</i> specifies a set of child processes for which status is requested. | DESCRIPTION | #include <sys wait.h=""> nid t waits id/nid t nid int % rat log int options).</sys> | SYNOPSIS #include <sys types.h=""></sys> | NAME waitpid – wait for child process to change state | waitpid(2) waitpid(2) |