accept(3) accept(3)

NAME

accept - accept a connection on a socket

SYNOPSIS

#include <sys/types.h>
#include <sys/socket.h>

int accept(int s, struct sockaddr *addr, int *addrlen);

DESCRIPTION

The argument s is a socket that has been created with socket(3N) and bound to an address with bind(3N), and that is listening for connections after a call to listen(3N). The accept() function extracts the first connection on the queue of pending connections, creates a new socket with the properties of s, and allocates a new file descriptor, ns, for the socket. If no pending connections are present on the queue and the socket is not marked as non-blocking, accept() blocks the caller until a connection is present. If the socket is marked as non-blocking and no pending connections are present on the queue, accept() returns an error as described below. The accept() function uses the accept() file to determine the STREAMS device file name associated with s. This is the device on which the connect indication will be accepted. The accepted socket, ns, is used to read and write data to and from the socket that connected to ns; it is not used to accept more connections. The original socket (s) remains open for accepting further connections.

The argument addr is a result parameter that is filled in with the address of the connecting entity as it is known to the communications layer. The exact format of the addr parameter is determined by the domain in which the communication occurs.

The argument *addrlen* is a value-result parameter. Initially, it contains the amount of space pointed to by *addr*; on return it contains the length in bytes of the address returned.

The accept() function is used with connection-based socket types, currently with SOCK_STREAM.

It is possible to **select**(3C) or **poll**(2) a socket for the purpose of an **accept**() by selecting or polling it for a read. However, this will only indicate when a connect indication is pending; it is still necessary to call **accept**().

RETURN VALUES

The accept() function returns -1 on error. If it succeeds, it returns a non-negative integer that is a descriptor for the accepted socket.

ERRORS

accept() will fail if:

EBADF The descriptor is invalid.

EINTR The accept attempt was interrupted by the delivery of a signal.

EMFILE The per-process descriptor table is full.

ENODEV The protocol family and type corresponding to s could not be found in the **netcon**

fig file.

ENOMEM There was insufficient user memory available to complete the operation.

EPROTO A protocol error has occurred; for example, the STREAMS protocol stack has not

been initialized or the connection has already been released.

EWOULDBLOCK The socket is marked as non-blocking and no connections are present to be

ccepted.

SEE ALSO

poll(2), bind(3N), connect(3N), listen(3N), select(3C), socket(3N), netconfig(4), attributes(5), socket(5)

bind(3)

NAME

bind - bind a name to a socket

SYNOPSIS

#include <sys/types.h>
#include <sys/socket.h>

int bind(int s, const struct sockaddr *name, int namelen);

DESCRIPTION

bind() assigns a name to an unnamed socket. When a socket is created with socket(3N), it exists in a name space (address family) but has no name assigned. bind() requests that the name pointed to by name be assigned to the socket.

RETURN VALUES

If the bind is successful, 0 is returned. A return value of -1 indicates an error, which is further specified in the global **errno**.

ERRORS

The **bind()** call will fail if:

EACCES The requested address is protected and the current user has inadequate permission

to access it.

EADDRINUSE The specified address is already in use.

EADDRNOTAVAIL The specified address is not available on the local machine.

EBADF s is not a valid descriptor.

EINVAL namelen is not the size of a valid address for the specified address family.

EINVAL The socket is already bound to an address.

ENOSR There were insufficient STREAMS resources for the operation to complete.

ENOTSOCK s is a descriptor for a fi le, not a socket.

The following errors are specific to binding names in the UNIX domain:

EACCES Search permission is denied for a component of the path prefix of the pathname in

name.

EIO An I/O error occurred while making the directory entry or allocating the inode.

EISDIR A null pathname was specified.

ELOOP Too many symbolic links were encountered in translating the pathname in *name*.

ENOENT A component of the path prefix of the pathname in *name* does not exist.

ENOTDIR A component of the path prefix of the pathname in *name* is not a directory.

EROFS The inode would reside on a read-only fi le system.

SEE ALSO

unlink(2), socket(3N), attributes(5), socket(5)

NOTES

Binding a name in the UNIX domain creates a socket in the file system that must be deleted by the caller when it is no longer needed (using **unlink**(2)).

The rules used in name binding vary between communication domains.

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fdopen(3) fdopen(3)

NAME

fdopen - associate a stream with a fi le descriptor

SYNOPSIS

#include <stdio.h>

FILE *fdopen(int fildes, const char *mode);

DESCRIPTION

The **fdopen()** function associates a stream with a fi le descriptor *fildes*, whose value must be less than 255.

The mode argument is a character string having one of the following values:

r or rb open a fi le for reading open a fi le for writing w or wh

a or ab open a fi le for writing at end of fi le open a fi le for update (reading and writing) r+ or rb+ or r+b w+ or wb+ or w+b open a fi le for update (reading and writing)

open a fi le for update (reading and writing) at end of fi le a+ or ab+ or a+b

The meaning of these flags is exactly as specified in **fopen**(3S), except that modes beginning with w do not cause truncation of the file.

The mode of the stream must be allowed by the file access mode of the open file. The file position indicator associated with the new stream is set to the position indicated by the file offset associated with the file

fdopen() will preserve the offset maximum previously set for the open file description corresponding to

The error and end-of-file indicators for the stream are cleared. The fdopen() function may cause the st atime fi eld of the underlying fi le to be marked for update.

RETURN VALUES

Upon successful completion, fdopen() returns a pointer to a stream. Otherwise, a null pointer is returned and errno is set to indicate the error.

fdopen() may fail and not set errno if there are no free stdio streams.

ERRORS

The fdopen() function may fail if:

EBADF The *fildes* argument is not a valid fi le descriptor.

EINVAL The *mode* argument is not a valid mode.

EMFILE FOPEN_MAX streams are currently open in the calling process.

EMFILE STREAM_MAX streams are currently open in the calling process.

ENOMEM Insuffi cient space to allocate a buffer.

USAGE

STREAM MAX is the number of streams that one process can have open at one time. If defi ned, it has the same value as FOPEN MAX.

File descriptors are obtained from calls like open(2), dup(2), creat(2) or pipe(2), which open files but do not return streams. Streams are necessary input for almost all of the Section 3S library routines.

SEE ALSO

creat(2), dup(2), open(2), pipe(2), fclose(3S), fopen(3S), attributes(5)

ip(7) ip(7)

NAME

ip - Linux IPv4 protocol implementation

#include <svs/socket.h>

SYNOPSIS

```
#include <netinet/in.h>
tcp_socket = socket(PF_INET, SOCK_STREAM, 0):
raw_socket = socket(PF_INET, SOCK_RAW, protocol);
udp socket = socket(PF INET, SOCK DGRAM, protocol);
```

DESCRIPTION

The programmer's interface is BSD sockets compatible. For more information on sockets, see **socket**(7).

An IP socket is created by calling the socket(2) function as socket(PF INET, socket type, protocol). Valid socket types are SOCK STREAM to open a tcp(7) socket, SOCK DGRAM to open a udp(7) socket, or SOCK_RAW to open a raw(7) socket to access the IP protocol directly. protocol is the IP protocol in the IP header to be received or sent. The only valid values for protocol are 0 and IPPROTO_TCP for TCP sockets and 0 and IPPROTO UDP for UDP sockets.

When a process wants to receive new incoming packets or connections, it should bind a socket to a local interface address using bind(2). Only one IP socket may be bound to any given local (address, port) pair. When INADDR ANY is specified in the bind call the socket will be bound to all local interfaces. When listen(2) or connect(2) are called on a unbound socket the socket is automatically bound to a random free port with the local address set to INADDR ANY.

ADDRESS FORMAT

An IP socket address is defined as a combination of an IP interface address and a port number. The basic IP protocol does not supply port numbers, they are implemented by higher level protocols like **tcp**(7).

```
struct sockaddr in {
    sa family t
                       sin family: /* address family: AF INET */
    u int16 t
                       sin port:
                                    /* port in network byte order */
    struct in_addr sin_addr;
                                     /* internet address */
/* Internet address. */
struct in addr {
    u int32 t
                       s addr;
                                     /* address in network byte order */
```

sin_family is always set to AF_INET. This is required; in Linux 2.2 most networking functions return **EINVAL** when this setting is missing. *sin port* contains the port in network byte order. The port numbers below 1024 are called reserved ports. Only processes with effective user id 0 or the **CAP_NET_BIND_SERVICE** capability may **bind**(2) to these sockets.

sin_addr is the IP host address. The addr member of struct in_addr contains the host interface address in network order. in_addr should be only accessed using the inet_aton(3), inet_addr(3), inet_makeaddr(3) library functions or directly with the name resolver (see **gethostbyname**(3)).

Note that the address and the port are always stored in network order. In particular, this means that you need to call htons(3) on the number that is assigned to a port. All address/port manipulation functions in the standard library work in network order.

SEE ALSO

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```
sendmsg(2), recvmsg(2), socket(7), netlink(7), tcp(7), udp(7), raw(7), ipfw(7)
```

sigaction(2) sigaction(2)

NAME

sigaction - POSIX signal handling functions.

SYNOPSIS

#include <signal.h>

int sigaction(int signum, const struct sigaction *act, struct sigaction *oldact);

DESCRIPTION

The sigaction system call is used to change the action taken by a process on receipt of a specific signal.

signum specifies the signal and can be any valid signal except SIGKILL and SIGSTOP.

If act is non-null, the new action for signal signum is installed from act. If oldact is non-null, the previous action is saved in oldact.

The sigaction structure is defi ned as something like

```
struct sigaction {
   void (*sa_handler)(int);
   void (*sa_sigaction)(int, siginfo_t *, void *);
   sigset_t sa_mask;
   int sa_flags;
   void (*sa_restorer)(void);
}
```

On some architectures a union is involved - do not assign to both sa_handler and sa_sigaction.

The sa_restorer element is obsolete and should not be used. POSIX does not specify a sa_restorer element.

sa_handler specifies the action to be associated with signum and may be SIG_DFL for the default action, SIG_IGN to ignore this signal, or a pointer to a signal handling function.

sa_mask gives a mask of signals which should be blocked during execution of the signal handler. In addition, the signal which triggered the handler will be blocked, unless the SA_NODEFER or SA_NOMASK flags are used.

sa_flags specifies a set of flags which modify the behaviour of the signal handling process. It is formed by the bitwise OR of zero or more of the following:

SA_NOCLDSTOP

If signum is SIGCHLD, do not receive notification when child processes stop (i.e., when child processes receive one of SIGSTOP, SIGTSTP, SIGTTIN or SIGTTOU).

SA_RESTART

Provide behaviour compatible with BSD signal semantics by making certain system calls restartable across signals.

RETURN VALUES

sigaction returns 0 on success and -1 on error.

ERRORS

EINVAL

An invalid signal was specified. This will also be generated if an attempt is made to change the action for **SIGKILL** or **SIGSTOP**, which cannot be caught.

SEE ALSO

```
kill(1), kill(2), killpg(2), pause(2), sigsetops(3),
```

sigsetops(3C) sigsetops(3C)

NAME

sigsetops, sigemptyset, sigfi llset, sigaddset, sigdelset, sigismember - manipulate sets of signals

SYNOPSIS

```
#include <signal.h>
int sigemptyset(sigset_t *set);
int sigfi llset(sigset_t *set);
int sigaddset(sigset_t *set, int signo);
int sigdelset(sigset_t *set, int signo);
int sigismember(sigset t *set, int signo);
```

DESCRIPTION

These functions manipulate sigset_t data types, representing the set of signals supported by the implementation

sigemptyset() initializes the set pointed to by set to exclude all signals defined by the system.

sigfi llset() initializes the set pointed to by set to include all signals defined by the system.

sigaddset() adds the individual signal specified by the value of signo to the set pointed to by set.

sigdelset() deletes the individual signal specified by the value of signo from the set pointed to by set.

sigismember() checks whether the signal specified by the value of *signo* is a member of the set pointed to by *set*.

Any object of type <code>sigset_t</code> must be initialized by applying either <code>sigemptyset()</code> or <code>sigfillset()</code> before applying any other operation.

RETURN VALUES

Upon successful completion, the **sigismember()** function returns a value of one if the specified signal is a member of the specified set, or a value of 0 if it is not. Upon successful completion, the other functions return a value of 0. Otherwise a value of -1 is returned and **errno** is set to indicate the error.

ERRORS

```
sigaddset(), sigdelset(), and sigismember() will fail if the following is true:
```

EINVAL The value of the *signo* argument is not a valid signal number.

sigfi llset() will fail if the following is true:

EFAULT The *set* argument specifi es an invalid address.

SEE ALSO

sigaction (2), sigpending (2), sigprocmask (2), sigsuspend (2), attributes (5), signal (5)

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socket(3) socket

NAME

socket - create an endpoint for communication

SYNOPSIS

```
cc [ flag ... ] fi le ... —lsocket —lnsl [ library ... ]
#include <sys/types.h>
#include <sys/socket.h>
```

int socket(int domain, int type, int protocol);

DESCRIPTION

socket() creates an endpoint for communication and returns a descriptor.

The *domain* parameter specifies a communications domain within which communication will take place; this selects the protocol family which should be used. The protocol family generally is the same as the address family for the addresses supplied in later operations on the socket. These families are defined in the include file <sys/socket.h>. There must be an entry in the netconfig(4) file for at least each protocol family and type required. If *protocol* has been specified, but no exact match for the tuplet family, type, protocol is found, then the first entry containing the specified family and type with zero for protocol will be used. The currently understood formats are:

PF_UNIX UNIX system internal protocols

PF INET ARPA Internet protocols

The socket has the indicated *type*, which specifies the communication semantics. Currently defined types are:

SOCK_STREAM SOCK_DGRAM SOCK_RAW SOCK_SEQPACKET SOCK RDM

A SOCK_STREAM type provides sequenced, reliable, two-way connection-based byte streams. An out-of-band data transmission mechanism may be supported. A SOCK_DGRAM socket supports datagrams (connectionless, unreliable messages of a fixed (typically small) maximum length). A SOCK_SEQPACKET socket may provide a sequenced, reliable, two-way connection-based data transmission path for datagrams of fixed maximum length; a consumer may be required to read an entire packet with each read system call. This facility is protocol specific, and presently not implemented for any protocol family. SOCK_RAW sockets provide access to internal network interfaces. The types SOCK_RAW, which is available only to the super-user, and SOCK_RDM, for which no implementation currently exists, are not described here.

protocol specifies a particular protocol to be used with the socket. Normally only a single protocol exists to support a particular socket type within a given protocol family. However, multiple protocols may exist, in which case a particular protocol must be specified in this manner. The protocol number to use is particular to the "communication domain" in which communication is to take place. If a protocol is specified by the caller, then it will be packaged into a socket level option request and sent to the underlying protocol layers.

Sockets of type SOCK_STREAM are full-duplex byte streams, similar to pipes. A stream socket must be in a connected state before any data may be sent or received on it. A connection to another socket is created with a connect(3N) call. Once connected, data may be transferred using read(2) and write(2) calls or some variant of the send(3N) and recv(3N) calls. When a session has been completed, a close(2) may be performed. Out-of-band data may also be transmitted as described on the send(3N) manual page and received as described on the recv(3N) manual page.

The communications protocols used to implement a SOCK_STREAM insure that data is not lost or duplicated. If a piece of data for which the peer protocol has buffer space cannot be successfully transmitted within a reasonable length of time, then the connection is considered broken and calls will indicate an error with -1 returns and with ETIMEDOUT as the specific code in the global variable errno. The protocols optionally keep sockets "warm" by forcing transmissions roughly every minute in the absence of other

socket(3) socket(3)

activity. An error is then indicated if no response can be elicited on an otherwise idle connection for a extended period (for instance 5 minutes). A **SIGPIPE** signal is raised if a process sends on a broken stream; this causes naive processes, which do not handle the signal, to exit.

SOCK_SEQPACKET sockets employ the same system calls as **SOCK_STREAM** sockets. The only difference is that **read**(2) calls will return only the amount of data requested, and any remaining in the arriving packet will be discarded.

SOCK_DGRAM and **SOCK_RAW** sockets allow datagrams to be sent to correspondents named in **sendto**(3N) calls. Datagrams are generally received with **recvfrom**(3N), which returns the next datagram with its return address.

An fcntl(2) call can be used to specify a process group to receive a SIGURG signal when the out-of-band data arrives. It may also enable non-blocking I/O and asynchronous notification of I/O events with SIGIO signals.

The operation of sockets is controlled by socket level *options*. These options are defined in the file <sys/socket.h>. setsockopt(3N) and getsockopt(3N) are used to set and get options, respectively.

RETURN VALUES

A -1 is returned if an error occurs. Otherwise the return value is a descriptor referencing the socket.

ERRORS

The **socket()** call fails if:

EACCES Permission to create a socket of the specified type and/or protocol is denied.

EMFILE The per-process descriptor table is full.

ENOMEM Insufficient user memory is available.

ENOSR There were insufficient STREAMS resources available to complete the opera-

tion.

EPROTONOSUPPORT The protocol type or the specified protocol is not supported within this

domain.

SEE ALSO

 $\label{eq:close} \begin{aligned} & \textbf{close}(2), \textbf{fcntl}(2), \textbf{ioctl}(2), \textbf{read}(2), \textbf{write}(2), \textbf{accept}(3N), \textbf{bind}(3N), \textbf{connect}(3N), \textbf{getsockname}(3N), \textbf{getsockopt}(3N), & \textbf{listen}(3N), & \textbf{recv}(3N), & \textbf{setsockopt}(3N), & \textbf{send}(3N), & \textbf{shutdown}(3N), & \textbf{socketpair}(3N), & \textbf{attributes}(5), \textbf{in}(5), \textbf{socket}(5) \end{aligned}$

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NAME

stat, fstat, lstat - get fi le status

SYNOPSIS

```
#include <sys/types.h>
#include <sys/stat.h>
#include <unistd.h>

int stat(const char * fi le_name, struct stat *buf);
int fstat(int fi ledes, struct stat *buf);
int lstat(const char * fi le name, struct stat *buf);
```

DESCRIPTION

These functions return information about the specified file. You do not need any access rights to the file to get this information but you need search rights to all directories named in the path leading to the file.

stat stats the fi le pointed to by fi le_name and fi lls in buf.

Istat is identical to **stat**, except in the case of a symbolic link, where the link itself is stat-ed, not the file that it refers to.

fstat is identical to **stat**, only the open fi le pointed to by *fi ledes* (as returned by **open**(2)) is stat-ed in place of *fi le name*.

They all return a stat structure, which contains the following fi elds:

```
struct stat {
 dev t
            st dev; /* device */
 ino_t
           st_ino; /* inode */
 mode_t st_mode; /* protection */
            st nlink; /* number of hard links */
 nlink t
 uid t
           st uid; /* user ID of owner */
 gid t
           st gid; /* group ID of owner */
 dev_t
           st_rdev; /* device type (if inode device) */
           st size: /* total size, in bytes */
 blksize t st blksize; /* blocksize for fi lesystem I/O */
 blkent t st blocks: /* number of blocks allocated */
            st atime: /* time of last access */
 time t
 time_t
            st mtime: /* time of last modification */
            st ctime; /* time of last status change */
 time t
```

The value *st_size* gives the size of the file (if it is a regular file or a symlink) in bytes. The size of a symlink is the length of the pathname it contains, without trailing NUL.

The value st_blocks gives the size of the file in 512-byte blocks. (This may be smaller than $st_size/512$ e.g. when the file has holes.) The value $st_blksize$ gives the "preferred" blocksize for efficient file system I/O. (Writing to a file in smaller chunks may cause an inefficient read-modify-rewrite.)

Not all of the Linux fi lesystems implement all of the time fi elds. Some fi le system types allow mounting in such a way that fi le accesses do not cause an update of the *st_atime* fi eld. (See 'noatime' in **mount**(8).)

The fi eld st_atime is changed by fi le accesses, e.g. by **execve**(2), **mknod**(2), **pipe**(2), **utime**(2) and **read**(2) (of more than zero bytes). Other routines, like **mmap**(2), may or may not update st_atime .

The fi eld st_mtime is changed by fi le modifications, e.g. by mknod(2), truncate(2), utime(2) and write(2) (of more than zero bytes). Moreover, st_mtime of a directory is changed by the creation or deletion of fi les in that directory. The st_mtime fi eld is not changed for changes in owner, group, hard link count, or mode.

stat(2)

The field st_ctime is changed by writing or by setting inode information (i.e., owner, group, link count, mode, etc.).

The following POSIX macros are defined to check the file type:

```
S_ISREG(m) is it a regular fi le?

S_ISDIR(m) directory?

S_ISCHR(m) character device?

S_ISBLK(m) block device?

S_ISFIFO(m) fi fo?

S_ISLNK(m) symbolic link? (Not in POSIX.1-1996.)

S_ISSOCK(m) socket? (Not in POSIX.1-1996.)
```

The following flags are defined for the *st_mode* field:

```
S IFMT
            0170000
                       bitmask for the fi le type bitfi elds
S IFSOCK
           0140000
S IFLNK
            0120000
                       symbolic link
S IFREG
            0100000
                       regular fi le
S IFBLK
            0060000
                       block device
S IFDIR
             0040000
                       directory
S IFCHR
            0020000
                       character device
S IFIFO
             0010000
                       fi fo
S ISUID
            0004000
                       set UID bit
S ISGID
            0002000
                       set GID bit (see below)
S ISVTX
            0001000
                       sticky bit (see below)
            00700
S IRWXU
                       mask for fi le owner permissions
S_IRUSR
            00400
                       owner has read permission
S IWUSR
            00200
                       owner has write permission
S IXUSR
            00100
                       owner has execute permission
S IRWXG
            00070
                       mask for group permissions
S_IRGRP
            00040
                       group has read permission
S IWGRP
            00020
                       group has write permission
S IXGRP
             00010
                       group has execute permission
S IRWXO
            00007
                       mask for permissions for others (not in group)
S IROTH
             00004
                       others have read permission
S IWOTH
            00002
                       others have write permisson
S IXOTH
            00001
                       others have execute permission
```

The set GID bit (S_ISGID) has several special uses: For a directory it indicates that BSD semantics is to be used for that directory; files created there inherit their group ID from the directory, not from the effective group ID of the creating process, and directories created there will also get the S_ISGID bit set. For a file that does not have the group execution bit (S_IXGRP) set, it indicates mandatory file/record locking.

The 'sticky' bit (S_ISVTX) on a directory means that a file in that directory can be renamed or deleted only by the owner of the file, by the owner of the directory, and by a privileged process.

RETURN VALUE

On success, zero is returned. On error, -1 is returned, and errno is set appropriately.

SEE ALSO

chmod(2), chown(2), readlink(2), utime(2), capabilities(7)

strerror(3) strerror(3)

NAME

strerror, strerror_r - return string describing error code

SYNOPSIS

#include <string.h>

```
char *strerror(int errnum);
int strerror_r(int errnum, char *buf, size_t n);
```

DESCRIPTION

The **strerror**() function returns a string describing the error code passed in the argument *errnum*, possibly using the LC_MESSAGES part of the current locale to select the appropriate language. This string must not be modified by the application, but may be modified by a subsequent call to **perror**() or **strerror**(). No library function will modify this string.

The **strerror_r**() function is similar to **strerror**(), but is thread safe. It returns the string in the user-supplied buffer buf of length n.

RETURN VALUE

The **strerror**() function returns the appropriate error description string, or an unknown error message if the error code is unknown. The value of *errno* is not changed for a successful call, and is set to a nonzero value upon error. The **strerror_r**() function returns 0 on success and -1 on failure, setting *errno*.

ERRORS

EINVAL

The value of errnum is not a valid error number.

ERANGE

Insuffi cient storage was supplied to contain the error description string.

CONFORMING TO

```
SVID 3, POSIX, BSD 4.3, ISO/IEC 9899:1990 (C89).
```

strerror_r() with prototype as given above is specified by SUSv3, and was in use under Digital Unix and HP Unix. An incompatible function, with prototype

```
char *strerror_r(int errnum, char *buf, size_t n);
```

is a GNU extension used by glibc (since 2.0), and must be regarded as obsolete in view of SUSv3. The GNU version may, but need not, use the user-supplied buffer. If it does, the result may be truncated in case the supplied buffer is too small. The result is always NUL-terminated.

SEE ALSO

```
errno(3), perror(3), strsignal(3)
```

waitpid(2) waitpid(2)

NAME

waitpid - wait for child process to change state

SYNOPSIS

```
#include <sys/types.h>
#include <sys/wait.h>
pid t waitpid(pid t pid, int *stat loc, int options);
```

DESCRIPTION

waitpid() suspends the calling process until one of its children changes state; if a child process changed state prior to the call to waitpid(), return is immediate. pid specifies a set of child processes for which status is requested.

If pid is equal to (**pid** t)-1, status is requested for any child process.

If pid is greater than $(pid_t)0$, it specifies the process ID of the child process for which status is requested.

If *pid* is equal to (**pid_t)0** status is requested for any child process whose process group ID is equal to that of the calling process.

If pid is less than (pid_t)-1, status is requested for any child process whose process group ID is equal to the absolute value of pid.

If **waitpid()** returns because the status of a child process is available, then that status may be evaluated with the macros defined by **wstat(5)**. If the calling process had specified a non-zero value of *stat_loc*, the status of the child process will be stored in the location pointed to by *stat_loc*.

The *options* argument is constructed from the bitwise inclusive OR of zero or more of the following flags, defined in the header <sys/wait.h>:

WCONTINUED The status of any continued child process specified by pid, whose status has not

been reported since it continued, is also reported to the calling process.

WNOHANG waitpid() will not suspend execution of the calling process if status is not imme-

diately available for one of the child processes specified by pid.

WNOWAIT Keep the process whose status is returned in stat_loc in a waitable state. The pro-

cess may be waited for again with identical results.

RETURN VALUES

If **waitpid()** returns because the status of a child process is available, this function returns a value equal to the process ID of the child process for which status is reported. If **waitpid()** returns due to the delivery of a signal to the calling process, **–1** is returned and **errno** is set to **EINTR**. If this function was invoked with **WNOHANG** set in *options*, it has at least one child process specified by *pid* for which status is not available, and status is not available for any process specified by *pid*, **0** is returned. Otherwise, **–1** is returned, and **errno** is set to indicate the error.

ERRORS

waitpid() will fail if one or more of the following is true:

ECHILD The process or process group specified by *pid* does not exist or is not a child of the call-

ing process or can never be in the states specified by options.

EINTR waitpid() was interrupted due to the receipt of a signal sent by the calling process.

EINVAL An invalid value was specified for options.

SEE ALSO

```
exec(2), exit(2), fork(2), sigaction(2), wstat(5)
```

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